

University of Puerto Rico Mayagüez
 Department of Electrical and Computer Engineering
 ICOM 4075: Foundations of Computing

Partial Exam I (100 Points)

Question 1 – (40 points) Short questions (2 points each)

- 1- Define the scope of a quantifier.

Part of an assertion in which variables are bound by the quantifiers.

$$\forall x (F(x) \wedge S(x)) \quad x: \text{wide scope}$$

$$\forall x F(x) \wedge \forall x S(x) \quad x: \text{narrow scope}$$

- 2- What is the truth set of the following predicate:

$$x^2 + 1 \leq 0 \text{ where the domain is } R$$

$$\emptyset$$

- 3- What is the inverse of the following implication?

I go to school only if it's not raining.

If I don't go to school then it is raining.

- 4- What is the converse of the following conditional proposition?

$$a^2 = b^2 \text{ unless } a \neq b$$

$$a^2 = b^2 \rightarrow a = b$$

- 5- What is the contrapositive of the following implication?

Snowing is sufficient for me to go fishing.

not

✓

If I don't go fishing then it's snowing.

- 6- If $P(x,y,z)$ is "z is an equation with x and y as its solutions" find two sets of domain for the following statement such that for one of them the statement is true and for the other one is false:

$$D_T = \{D_z, D_{x,y}\} \quad D_F = \{D'_z, D'_{x,y}\}$$

$$\exists z \forall x \forall y \neg p(x,y,z)$$

$$D_z = \text{Set of all equations on } \mathbb{R} \quad D_{x,y} = \mathbb{R} \quad D'_z = \{x+y=0\} \quad D'_{x,y} = \mathbb{Z}$$

- 7- Translate the following statement into the logical expression (using quantifiers and predicates) in two ways; one with the domain as people and the other with the domain as students:

$C(x)$: x is in class Some students in your class do not want to be rich Domain students: $\exists x C(x) \wedge \neg R(x)$

$R(x)$: x wants to be rich

$S(x)$: x is student $SC(x)$: x is student in class

Domain People \leftarrow Answer 1: $\exists x SC(x) \wedge \neg R(x)$

Answer 2: $\exists x S(x) \wedge C(x) \wedge \neg R(x)$

- 8- Negate the following proposition such that no negate sign comes before a quantifier:

$$\forall x \exists y (p(x,y) \rightarrow \forall z (q(x,z,10) \oplus q(x,z,11)))$$

$$\exists x \forall y (\neg p(x,y) \wedge \exists z \neg (\neg q(x,z,10) \oplus q(x,z,11)))$$

9- Parenthesize: $(p \rightarrow (q \vee r)) \rightarrow ((\neg s \wedge p) \rightarrow (\neg q \wedge r))$

Translate the statements of questions 10 to 16 assuming that $L(x,y)$ is "x loves y" and $H(x,y)$ is "x hates y" and the domain for both x and y are all people.

10- $\exists x \exists y \exists z (x \neq y) \wedge L(z,x) \wedge H(z,y)$

There are two different persons x, y and a person z. That z loves x and hates y.

11- $\exists x L(x, Juan) \wedge H(x, Hector)$

There is a person who loves Juan and hates Hector.

12- $\exists x L(x, Juan) \wedge (\forall z L(z, Juan) \rightarrow z = x)$

There is a unique person who loves Juan.

13- $L(Mary, Juan) \wedge L(Sarah, Juan) \rightarrow H(Mary, Sarah) \vee H(Sarah, Mary)$

If both Mary and Sarah love Juan, either Mary hates Sarah or Sarah hates Mary.

14- $\forall x \exists y \forall z L(x, z) \rightarrow H(y, z)$

For any person x, there are some people y such that whoever x loves, y hates.

15- There is one and only one person that if you love him he loves you too.

$$\exists ! x \quad L(\text{you}, x) \rightarrow L(x, \text{you})$$

16- There is at least one person who if Juan doesn't love her, she would hate Juan.

$$\exists x \quad \neg L(Juan, x) \rightarrow H(x, Juan)$$

Multiple choice questions from 17 to 20:

17- Which one is not tautology?

- a) $\forall x P(x) \rightarrow P(x) \vee Q(x)$
- b)** $(p \rightarrow q) \wedge (p \rightarrow \neg q)$
- c) $P(x)$ for arbitrary $x \rightarrow \forall x P(x)$
- d) $\neg \exists x P(x) \leftrightarrow \forall x \neg P(x)$

18- In an exam one question was: "prove that there are no solutions in positive integers for the equation $x^3 + y^3 = 21$ ". One student wrote the following: "if $x^3 + y^3 = 21$ since $3^3=27$ then x and y must be less than 3. Thus the possibilities for x and y are 1 or 2. But none of $1^3+2^3, 1^3+1^3, 2^3+2^3$ makes 21".

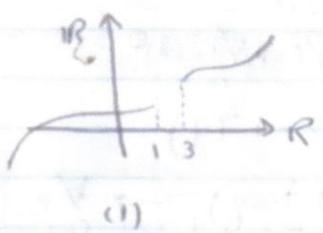
- a) This is a proof by example
- b) This is proof by cases
- c) This is a disproof by example
- d) This is a wrong proof

- ~~excluded~~ 19- Which one is not a tautology?
- a) $(p \rightarrow q) \leftrightarrow (q \vee \neg p)$
 - b) $p \wedge q \rightarrow p \vee q$
 - c) $\exists x P(x) \rightarrow P(a)$ for $a \in U$ where U is the domain
 - d) $\exists x (P(x) \vee Q(x)) \leftrightarrow \exists x P(x) \vee \exists x Q(x)$

20- Which one is a function (Place an X next to it):

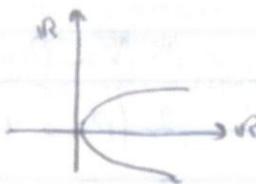
20- which one is a function

$$f: \mathbb{R} \rightarrow \mathbb{R}$$



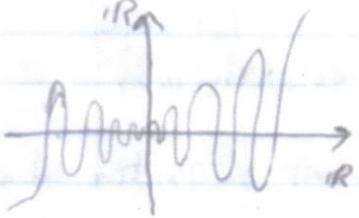
(1)

$$f: \mathbb{R}^+ \rightarrow \mathbb{R}$$



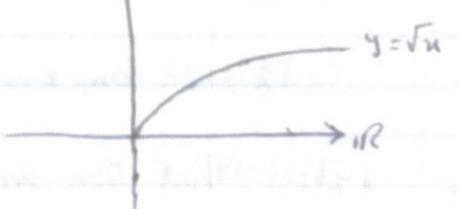
(2)

$$f: \mathbb{R} \rightarrow \mathbb{R}$$



(3)

$$f: \mathbb{R} \rightarrow \mathbb{R}$$



(4)

- 1: $\exists x \neg P(x)$ Premise
 2: $\neg P(a)$ EI
 3: $\forall x P(x) \vee Q(x)$ Premise
 4: $P(a) \vee Q(a)$ UI (in Particular for a)

Question 2 (10 points) Show that the following argument is valid:

$$\begin{array}{l}
 \forall x P(x) \vee Q(x) \\
 \forall x \neg Q(x) \vee S(x) \\
 \forall y R(y) \rightarrow \neg S(y) \\
 \exists x \neg P(x) \\
 \hline
 \therefore \exists x \neg R(x)
 \end{array}$$

- 5: $Q(a)$
 6: $\forall x \neg Q(x) \vee S(x)$ Premise
 7: $\neg Q(a) \vee S(a)$ UI (in Particular for a)
 8: $S(a)$
 9: $\forall y R(y) \rightarrow \neg S(y)$ Premise
 10: $R(a) \rightarrow \neg S(a)$ UI (in Particular for a)
 11: $S(a) \rightarrow \neg R(a)$ contrapositive of 10
 12: $\neg R(a)$
 13: $\exists x \neg R(x)$ EG

together with inference rules,
 equivalences are allowed also.
 Can be inferred directly
 from 8 & 10 & MT

Question 3 (10 points) Provide a proof for the following two equivalences using the specified method:

- a) $(p \rightarrow q) \wedge \neg p \equiv p \rightarrow (q \wedge \neg p)$ Not using truth table
 b) $(p \rightarrow q) \oplus p \equiv p \rightarrow (q \oplus p)$ Using truth table

a) $(p \rightarrow q) \wedge \neg p \equiv (q \vee \neg p) \wedge \neg p \equiv (q \wedge \neg p) \vee (\neg p \wedge \neg p) \equiv (q \wedge \neg p) \vee \neg p \equiv p \rightarrow (q \wedge \neg p)$

b)

P	q	$p \rightarrow q$	$q \oplus p$	$(p \rightarrow q) \oplus p$	$p \rightarrow (q \oplus p)$
T	T	T	F	F	F
T	F	F	T	T	T
F	T	T	T	T	T
F	F	T	F	T	T

Question 4-(10 points) Prove that any 2-digit integer number ab satisfies the following property:

$$ab - ba = 9k$$

(Note: for example in a 2-digit number like 64, $a=6$ and $b=4$. Hint: Since the number is in base 10, try to use the representation of numbers in base 10)

What kind of proof did you use?

$$ab = b + 10a$$

$$ba = a + 10b$$

$$ab - ba = b + 10a - a - 10b = 9a - 9b = 9(a - b)$$

Direct proof

Question 5-(10 points) If $f: A \rightarrow B$ and $g: B \rightarrow C$ are functions and f and g are 1-1 and onto, show that

$g \circ f$ is a bijection.

$$g \circ f(x_1) = g(f(x_1)) \underset{f}{\underset{\sim}{\rightarrow}} g(f(x_2)) \underset{g}{\underset{\sim}{\rightarrow}} g(y_1) = g(y_2)$$

$$(g \text{ is 1-1}) \wedge (g(y_1) = g(y_2)) \rightarrow y_1 = y_2 \rightarrow f(x_1) = f(x_2)$$

$$(f \text{ is 1-1}) \text{ thus } (f(x_1) = f(x_2) \rightarrow x_1 = x_2)$$

Take an arbitrary element $c \in C$. Since g is onto, there exist an element $b \in B$ s.t. $g(b) = c$. For $b \in B$ there exist $a \in A$ s.t.

$f(a) = b$ because f is onto.

$$g \circ f(a) = g(f(a)) = g(b) = c$$

That means for any $c \in C$ there exist an $a \in A$ s.t.

$g \circ f(a) = c$. Therefore $g \circ f$ is onto.

Question 5- (10 points) Proof the followings set propositions:

- a) $A \times B = \emptyset \rightarrow A = \emptyset \vee B = \emptyset$ Using contradiction
- b) $(A \cup B) - (A \cap B) = (A - B) \cup (B - A)$ Using 2-way subsets
- c) $(A \cap B) \cup (A \cap B^c) = A$ Using identity laws

$$a) \neg(A = \emptyset \vee B = \emptyset) \equiv A \neq \emptyset \wedge B \neq \emptyset$$

$$\begin{aligned} A \neq \emptyset &\rightarrow \exists a \in A \\ B \neq \emptyset &\rightarrow \exists b \in B \end{aligned} \quad \left. \begin{array}{l} \rightarrow (a, b) \in A \times B \\ \text{by The definition of } A \times B \end{array} \right.$$

$$\rightarrow A \times B \neq \emptyset \equiv \neg(A \times B = \emptyset)$$

$$b) x \in (A \cup B) - (A \cap B) \rightarrow x \in A \cup B \wedge x \notin A \cap B \rightarrow x \in A \cup B \wedge x \in \overline{A \cap B}$$

$$\rightarrow x \in A \cup B \wedge x \in \overline{A} \cup \overline{B} \rightarrow (x \in A \vee x \in B) \wedge (x \in \overline{A} \vee x \in \overline{B})$$

$$\rightarrow ((x \in A \vee x \in B) \wedge x \in \overline{A}) \vee ((x \in A \vee x \in B) \wedge x \in \overline{B})$$

$$\rightarrow (x \in A \wedge x \in \overline{A}) \vee (x \in B \wedge x \in \overline{A}) \vee (x \in A \wedge x \in \overline{B}) \vee (x \in B \wedge x \in \overline{B})$$

$$\rightarrow x \in A \cap \overline{A} \vee x \in B \cap \overline{A} \vee x \in A \cap \overline{B} \vee x \in B \cap \overline{B} \rightarrow F \vee x \in B - A \vee x \in A - B \vee F$$

$$\rightarrow x \in B - A \vee x \in A - B \rightarrow x \in (B - A) \cup (A - B)$$

$$\text{Therefore } (A \cup B) - (A \cap B) \subseteq (B - A) \cup (A - B)$$

The other side:

$$x \in (B - A) \cup (A - B) \rightarrow x \in B - A \vee x \in A - B \rightarrow (x \in B \wedge x \notin A) \vee (x \in A \wedge x \notin B)$$

$$\rightarrow ((x \in B \wedge x \notin A) \vee x \in A) \wedge ((x \in B \wedge x \notin A) \vee x \notin B) \rightarrow (x \in B \vee x \in A) \wedge (x \notin A \vee x \notin B)$$

$$(x \in B \vee x \notin B) \wedge (x \notin A \vee x \notin B) \rightarrow x \in B \cup A \wedge (x \in \overline{A} \vee x \in A) \wedge (x \in \overline{B} \vee x \in B) \wedge (x \in \overline{A} \vee x \in \overline{B})$$

$$\rightarrow x \in B \cup A \wedge x \in A \cap \overline{A} \wedge x \in B \cap \overline{B} \wedge x \in \overline{A} \cup \overline{B} \rightarrow x \in B \cup A \wedge T \wedge T \wedge x \in \overline{A} \cup \overline{B} \rightarrow$$

$$x \in B \cup A \wedge x \in \overline{A \cap B} \rightarrow x \in B \cup A \wedge x \notin \overline{A \cap B} \rightarrow x \in (B \cup A) - (A \cap B)$$

$$\text{Therefore } (B - A) \cup (A - B) \subseteq (A \cup B) - (A \cap B)$$

$$c) (A \cap B) \cup (A \cap \overline{B}) = A \cap (B \cup \overline{B}) = A \cap U = A$$

Question 7- (10 points) Show that if $x \in R$ and $n \in Z$ then $x < n \leftrightarrow [x] < n$

we have the property

$$[x] \leq x < [x]+1$$

assume $x < n$

$$[x] \leq x \wedge x < n \rightarrow [x] < n$$

The other side : assume $[x] < n$

$$[x] < n \rightarrow [x] \leq n-1$$

$$x < [x]+1 \wedge [x] \leq n-1 \rightarrow x < n-1+1$$

$$\rightarrow x < n$$

Question 8- (BONUS 10 points)- If $f: A \rightarrow B$ is a function and $S \subseteq A, T \subseteq A$ and $L \subseteq B$ show that:

a) $f(S \cap T) \subseteq f(S) \cap f(T)$

b) $f^{-1}(L^c) = (f^{-1}(L))^c$

a) $y \in f(S \cap T) \rightarrow \exists x \in S \cap T, f(x) = y$

$$\rightarrow (x \in S \wedge f(x) = y) \wedge (x \in T \wedge f(x) = y)$$

$$\rightarrow y \in f(S) \wedge y \in f(T) \rightarrow y \in f(S) \cap f(T)$$

b) $f^{-1}(\bar{L}) = \{x \in A \mid f(x) \in \bar{L}\} = \{x \in A \mid f(x) \notin L\}$

$$= \{x \in A \mid x \notin f^{-1}(L)\} = \{x \in A \mid x \in \bar{f^{-1}(L)}\}$$

Note: for part a, it is not necessary for f to be invertible. But if you give a proof in which you assume that f is invertible, we accept it.